

## Solutions

1. (a) Any part of an optimal path is itself optimal B1
- (b) The route chosen such that the maximum arc length is as small as possible B1
- (c) e.g. Maximising freight by minimising fuel needed when planning multiple stage light aircraft journey B2, 1, 0  
*B1 cao ("port", "section", OK; "arc", "stage", activity", "event", not)*  
*B1 cao (not min of max rate, not minimize largest arc)*  
*B2 cao*  
*B1 cloze "Bod" gets B1*

[4]

2. Let  $x_{ij} = 1$  if worker does task, 0 otherwise B1  
 where  $x_{ij}$  indicates the arc from node  $i$  to node  $j$  i.e P, Q, R  $j \in \{1, 2, 3\}$  B1

$$\begin{array}{llll} x_{p1} + x_{p2} + x_{p3} = 1 & x_{p1} + x_{q1} + x_{r1} = 1 & M1 & \\ x_{q1} + x_{q2} + x_{q3} = 1 & \text{and } x_{p2} + x_{q2} + x_{r2} = 1 & A1 & \\ x_{r1} + x_{r2} + x_{r3} = 1 & x_{p3} + x_{q3} + x_{r3} = 1 & A1 & 3 \end{array}$$

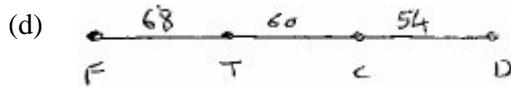
Minimise,  $C = 8x_{p1} + 7x_{p2} + 3x_{p3} + 9x_{q1} + 5x_{q2} + 6x_{q3} + 10x_{r1} + 4x_{r2} + 4x_{r3}$   
 where  $C$  is in hundreds of pounds B1, B1 2

*B1 cao*  
*B1 defining variable – attempt*  
*M1 at least 3 equations – coefficients of one*  
*A1 cao 3 correct*  
*A1 cao 6 correct*  
*B1 Minimise*  
*B1 cao (condone a slip) (- accept cost in pounds)*

[7]

3. (a) Each activity must be visited once and then we return to the starting activity, this must be done in a minimum time B2, 1, 0 2  
*B2 cao – all 3 bits in the context*  
*B1 cloze 'Bod' is B1 (e.g. not in context; just 'each activity once' – but not all 3; ...)*
- (b)  $108 + 54 + 150 + 68 + 100 = 480$  minutes (= 8 hours) M1 A1 2  
*M1 (maybe implicit) attempting to add 5 values*  
*A1 cao*

- (c) Use nearest neighbour B F T C D B M1 A1  
 $64 + 68 + 60 + 54 + 150 = 396$  minutes (67 hours) A1 3  
*M1 each vertex visited once – either NN or 2 x mst-shortcut (BD)*  
*A1 cao incl return to B (BFTCDB)*  
*A1 cao (396)*



- CT, TF, CD (Prim or Kruskal) M1 A1  
 $182 + 64 + 100 = 346$  minutes M1 A1ft 4  
*M1 Finding correct minimum spanning tree (maybe implicit) 182 sufficient*  
*A1 cao tree or 182*  
*M1 adding 2 least arcs to B i.e. 100 and 64 only*  
*A1ft cao ft from their m.s.t. value i.e. 164 and their tree length*

[11]

4. (a) Adding  $n \geq 20$  to table to give B1

	H	P	R	W
A	3	5	11	9
B	3	7	8	N
C	2	5	10	7
D	8	3	7	6

Reducing rows first  $\begin{bmatrix} 0 & 2 & 8 & 6 \\ 0 & 4 & 5 & n-3 \\ 0 & 3 & 8 & 5 \\ 5 & 0 & 4 & 3 \end{bmatrix}$  then columns  $\begin{bmatrix} 0 & 2 & 4 & 3 \\ 0 & 4 & 1 & n-6 \\ 0 & 3 & 4 & 2 \\ 5 & 0 & 0 & 0 \end{bmatrix}$  M1 A13

Either  $\begin{bmatrix} 0 & 1 & 3 & 2 \\ 0 & 3 & 0 & n-7 \\ 0 & 2 & 3 & 1 \\ 6 & 0 & 0 & 0 \end{bmatrix}$  or  $\begin{bmatrix} 0 & 1 & 3 & 2 \\ 0 & 3 & 0 & n-7 \\ 0 & 2 & 3 & 1 \\ 6 & 0 & 0 & 0 \end{bmatrix}$  M1 A1ft

↓ ↓

$\begin{bmatrix} 0 & 0 & 2 & 1 \\ 1 & 3 & 0 & n-7 \\ 0 & 1 & 2 & 0 \\ 7 & 0 & 0 & 0 \end{bmatrix}$   $\begin{bmatrix} 0 & 0 & 3 & 1 \\ 0 & 2 & 0 & n-8 \\ 0 & 1 & 3 & 0 \\ 7 & 0 & 1 & 0 \end{bmatrix}$  M1 A1ft 4

- A – H P  
 B – R or R cost £21 000  
 C – W H A1

D – P W

A1 2

(b) Not unique – gives the other solution

M1 A1ft 2

[11]

5.

Stage	State	Action	Value
1	H	HT	4*
	I	IT	3*
	J	JT	12*
	K	KT	20*
2	D	DH	2 + 4 = 6
		DI	4 + 3 = 7*
	E	EH	3 + 4 = 7*
		EI	4 + 3 = 7*
	F	FJ	10 + 12 = 22*
		FK	-8 + 20 = 12
	G	GJ	10 + 12 = 22
		GK	17 + 20 = 37*
3	A	AD	3 + 7 = 10
		AE	2 + 7 = 9
		AF	-5 + 22 = 17*
	B	BD	3 + 7 = 10
		BE	2 + 7 = 9
		BF	-6 + 22 = 16*
	C	CF	8 + 22 = 30*
		CG	-15 + 37 = 22
		SA	2 + 17 = 19
4	S	SB	3 + 16 = 19
		SC	-10 + 30 = 20*

M1 A1 2

M1 A1

A1 3

M1 A1ft

A1 ft 3

M1 A1ft 2

Route S C F J T £20 000

M1 A1 2

[12]

6. (a) Either e.g.

In an  $n \times m$  problem, a degenerate solution occurs when the number of cells used is less than  $(n + m - 1)$

B2,1,0 2

or e.g. when all the demand for one destination is satisfied by all the supply from a source, before the final demand and supplies are allocated

*B2 cao*

*B1 cloze "bod" is B1*

- (b) If the total supply > total demand a dummy is used to absorb the excess  
*B1 cao must (cannot decipher copy properly)*

B1 1

(c) 
$$\begin{bmatrix} 15 & & \\ 1 & 11 & 0 \\ & & 17 \end{bmatrix}$$

B1 1

*B1 cao total of five numbers*

- (d) Shadow costs  $S_A = 0 \quad S_B = -1 \quad S_C = -1$   
 $D_1 = 62 \quad D_2 = 49 \quad D_3 = 1$

Improvement indices  $I_{A2} = 47 - 0 - 49 = -2^*$   
 $I_{A3} = 0 - 0 - 1 = -1$   
 $I_{C1} = 68 + 1 - 62 = 7$   
 $I_{C2} = 58 + 1 - 49 = 10$

	1 <sup>(62)</sup>	2 <sup>(49)</sup>	3 <sup>(1)</sup>
⊕ A	15-θ	θ	
⊖ B	1+θ	11-θ	0
⊖ C			17

M1A1A1ft 3

*Entering A2, exiting B2, θ = 0*

- Shadow costs  $S_A = 0 \quad S_B = -1 \quad S_C = -1$   
 $D_1 = 62 \quad D_2 = 47 \quad D_3 = 1$

Improvement indices  $I_{A3} = 0 - 0 - 1 = -1^*$   
 $I_{B2} = 48 + 1 - 47 = 2$   
 $I_{C1} = 68 + 1 - 62 = 7$   
 $I_{C2} = 58 + 1 - 47 = 12$

	1 <sup>(62)</sup>	2 <sup>(47)</sup>	3 <sup>(1)</sup>
⊕ A	4-θ	11	θ
⊖ B	12+θ		0-θ
⊖ C			17

M1A1A1ft 3

*Entering A3, exiting B3, θ = 0*

	1 <sup>(62)</sup>	2 <sup>(47)</sup>	3 <sup>(1)</sup>
⊕ A	4	11	θ
⊖ B	12		
⊖ C			17

- Shadow costs  $S_A = 0 \quad S_B = -1 \quad S_C = 0$   
 $D_1 = 62 \quad D_2 = 47 \quad D_3 = 0$

M1 A1

Improvement indices  $I_{B2} = 48 + 1 - 47 = 2$   
 $I_{B3} = 0 + 1 - 0 = 1$   
 $I_{C1} = 68 - 0 - 62 = 6$  B1  
 $I_{C2} = 58 - 0 - 47 = 11$

∴ Optimal

Cost 1497 units B1 4

[14]

7. (a) e.g. Maximise  $P = V$  B1  
 Subject to:  $V - 5p_1 - 3p_2 - 6p_3 + r = 0$  M1  
 $V - 7p_1 - 8p_2 - 4p_3 + s = 0$  A2,1,0  
 $V - 2p_1 - 4p_2 - 9p_3 + t = 0$   
 $p_1 + p_2 + p_3 (+u) = 1$

where  $V$  = value of game to A,  $P_i$  = probability of A playing row  $i$   
 $P_i \geq 0$  and  $r, s, t, u$  are slack variables all  $\geq 0$

B1 5

*B1 Maximise/minimise and consistent function*  
*M1 constraints (condone non-negativity)*  
*– at least one correct must be equations*  
*A2 all correct*  
*A1 at least two correct*  
*B1 defining variables*

- (b) Not reducible and a three variable problem B1 1  
*B1 cao – both*

(c) e.g.

b v	V	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	r	s	t	u	value	
r	1	-5	-3	-6	1	0	0	0	0	M1
s	1	-7	-8	-4	0	1	0	0	0	A1
t	1	-2	-4	-9	0	0	1	0	0	2
u	0	1	1	1	0	0	0	1	1	
P	-1	0	0	0	0	0	0	0	0	

b v	V	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	r	s	t	u	value	Row ops
V	1	-5	-3	-6	1	0	0	0	0	R <sub>1</sub> / 1 M1 A1
s	0	-2	-5	-4	-1	1	0	0	0	R <sub>2</sub> - R <sub>1</sub> A1
t	0	-3	-1	-3	-1	0	1	0	0	R <sub>3</sub> - R <sub>1</sub> B1ft
u	0	1	1	1	0	0	0	1	1	R <sub>4</sub> stet 4
P	0	-5	-3	-6	1	0	0	0	0	R <sub>5</sub> + R <sub>1</sub>

b v	V	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	r	s	t	u	value	Row ops
V	1	-11	-18	0	-2	3	0	0	0	R <sub>1</sub> + 6R <sub>2</sub> M1 A1ft

$P_3$	0	-1	$-\frac{5}{2}$	1	$\frac{1}{2}$	$\frac{1}{2}$	0	0	0	$R_2/2$	A1
t	0	0	$-\frac{17}{2}$	0	$-\frac{5}{2}$	$\frac{5}{2}$	1	0	0	$R_3 + 3R_2$	B1ft
u	0	2	$\frac{7}{2}$	0	$\frac{1}{2}$	$-\frac{1}{2}$	0	1	1	$R_4 - R_2$	4
P	0	-11	-18	0	-2	3	0	0	0	$R_5 + 6R_2$	

[16]

8. (a)  $7x + 10y + 10z + r = 3600$

$6x + 9y + 12z + s = 3600$

B2,1,0

$2x + 3y + 4z + t = 2400$

$P - 35x - 55y - 60z = 0$

B2,0 4

(b) (i)

b.v.	x	y	z	r	s	t	value	Row ops	
r	2	$\frac{5}{2}$	0	1	$-\frac{5}{6}$	0	600	$R_1 - 10R_2$	A1
z	$\frac{1}{2}$	$\frac{3}{4}$	1	0	$\frac{1}{12}$	0	300	$R_2 \div 12$	M1
t	0	0	0	0	$-\frac{1}{3}$	1	1200	$R_3 - 4R_2$	A1ft
P	-5	-10	0	0	5	0	1800	$R_4 + 60R_2$	B1 5

(ii)

b.v.	x	y	z	r	s	t	value	Row ops	
y	$\frac{4}{5}$	1	0	$\frac{2}{5}$	$-\frac{1}{3}$	0	240	$R_1 \div \frac{5}{2}$	M1
z	$-\frac{1}{10}$	0	1	$-\frac{3}{10}$	$\frac{1}{3}$	0	120	$R_2 - \frac{3}{4}R_1$	A1ft
t	0	0	0	0	$-\frac{1}{3}$	1	1200	$R_3$ stet	M1
P	3	0	0	4	$\frac{5}{3}$	0	20400	$R_4 + 10R_1$	A1 4

(c)  $P = 20400$

$x = 0$

$y = 240$

$z = 120$

M1

$r = 0$

$s = 0$

$t = 1200$

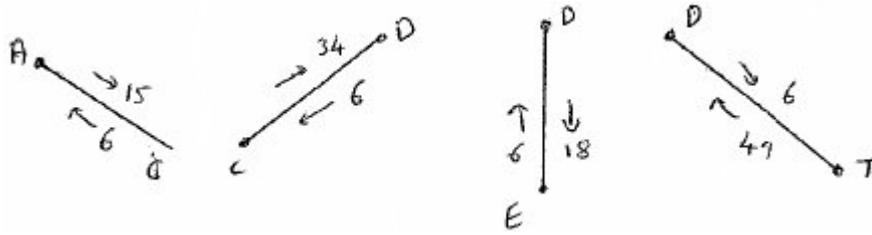
A2ft, A1ft, 0

2

[16]

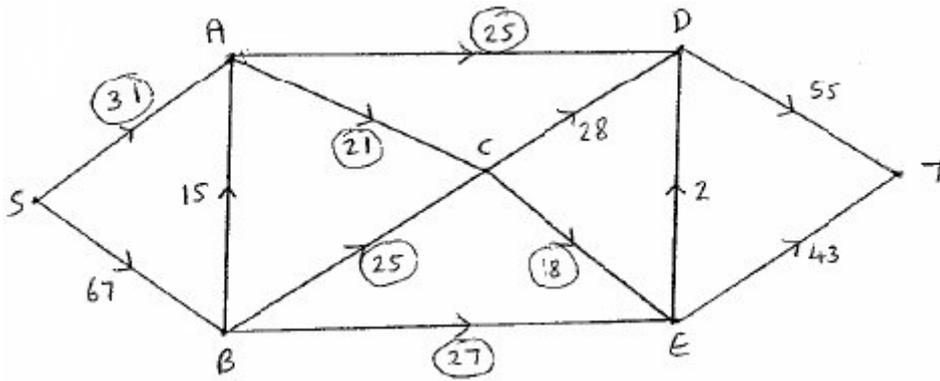
9. (a)  $C_1 = 103$ ,  $C_2 = 177$ , flow = 76 B1, B1, B1 3

(b) M1A1 2



(c) e.g. SBCDT - 6 M1  
 SBCDET - 1 A3,2,1,0  
 SBACDET - 15 B1 5  
 Max flow is 98

(d) M1A1 2



(e) Maximum flow = minimum cut M1  
 Cut through AD, AC, BC and BE A1 2